# Improving AMG interpolation through energy minimization

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#### 1 Motivation

- Energy minimimaztion
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AMG is a very popular solution algorithm which is used in several applications:

- it is very flexible as it only requires little information beyond the matrix itself;
- it can be parallelized up to millions of cores;
- as preconditioner for a Krylov method often guarantees performance independent of the grid size;

AMG works almost perfectly for Poisson problems, however, it is less effective in more difficult problems such as those arising in structural mechanics:

- classical or extended interpolation usually give poor results;
- aggregation-based AMG is generally preferred over Classical AMG;
- the near kernel space is larger than the standard constant vector (which is fine for Poisson);



In the context of structural problems (or other difficult problems):

- more powerful smoothers than Jacobi or Gauss-Seidel are needed;
- the near kernel is constructed starting from the Rigid Body Modes (RBM);
- it is often necessary to improve the near kernel through smoothing or an eigensolution;
- prolongation is constructed through a least square fit of the test space;

Main issues:

- coarse basis functions with local support are not orthogonal to high frequency modes;
- the prolongation has to be smoothed to allow for fast convergence (helpful in calssical AMG as well);
- smoothed prolongation however leads to high complexities;



The main idea is, for a given prolongation pattern  $\mathcal{P}$ , to compute:  $P = \underset{P \in \mathcal{P}}{\operatorname{argmin}} \left( \operatorname{tr}(P^T A P) \right)$ 

or, equivalently, to minimize the energy of every prolongation column:

 $p_i^T A p_i \to \min \quad \forall i \in C$ 

We consider a **Classical** AMG framework, so that:

$$A = \begin{bmatrix} A_{ff} & A_{fc} \\ A_{fc}^T & A_{cc} \end{bmatrix} \qquad P = \begin{bmatrix} W \\ I \end{bmatrix} \qquad R = \begin{bmatrix} 0 & I \end{bmatrix} \qquad S = \begin{bmatrix} I \\ 0 \end{bmatrix}$$

 Olson, L. N., Schroder, J. B. & Tuminaro, R. S., A General Interpolation Strategy for Algebraic Multigrid Using Energy Minimization, SIAM Journal on Scientific Computing 33, pp. 966–991 (2011).

2 Manteuffel, T. A., Olson, L. N., Schroder, J. B. & Southworth, B. S., A Root-Node-Based Algebraic Multigrid Method, SIAM Journal on Scientific Computing 39, S723–S756 (2017).



Suppose we have a C/F partition of the unknowns and an SPD X matrix which is spectrally equivalent to the symmetrized smoother  $\widetilde{M} = M^T (M + M^T - A)^{-1} M$ .

Defining 
$$A_s = S^T A S$$
 and  $X_s = S^T X S$  and:  
 $k_S \le \lambda_{min}(X_S^{-1} A_S) \le \lambda_{max}(X_S^{-1} A_S) \le c_2$ 

then, if  $||PR||_A$  is bounded:

$$K_{TG} \leq \frac{c_2}{k_S} \|PR\|_A^2$$
 where  $K_{TG}$  is such that  $\|E_{TG}\|_A \leq 1 - \frac{1}{K_{TG}}$ 

In practice, the smaller  $||PR||_A$ , the faster convergence!

Falgout R.D. and Vassilevski P. S., On Generalizing the Algebraic Multigrid Framework, SIAM Journal on Numerical Analysis 42, pp. 1669–1693 (2004).



The minimum in energy can be simply found by deriving with respect to W and setting to zero:

$$\frac{\partial}{\partial W} \left( \operatorname{tr}(P^T A P) \right) = 0 \Longrightarrow A_{ff} W = -A_{fc}$$

which is exactly what is defined *ideal* prolongation  $W_{id} = -A_{ff}^{-1}A_{fc}$ .

Note that it is also the target that is sought with a 2x2 block FSAI approximation:

$$A = \begin{bmatrix} A_{11} & A_{12} \\ A_{12}^T & A_{22} \end{bmatrix}, \quad G = \begin{bmatrix} I & 0 \\ F & I \end{bmatrix}, \quad GAG^T \simeq I \qquad \Rightarrow \qquad F = -A_{11}A_{12}$$

Janna, C., Ferronato, M. & Gambolati, G., A Block FSAI-ILU Parallel Preconditioner for Symmetric Positive Definite Linear Systems, SIAM Journal on Scientific Computing 32, pp. 2468–2484 (2010).



$$\overline{P} = S^k \overline{P}_0$$

for a small power k.

Then we minimize the energy of every column of P restricted to the non-zero pattern  $\overline{P}$  by solving the following sequence of  $n_c$  dense linear systems:

$$A(I_i, I_i)\widetilde{p}_i = -A(I_i, i) \quad \forall i \in C$$

with  $I_i$  the set of indices of the non-zeroes in the *i*-th column of  $\overline{P}$ .



$$V \subseteq \mathsf{range}(P), \qquad \mathsf{with} \qquad V = \begin{bmatrix} V_f \\ V_c \end{bmatrix}$$

This condition becomes:

$$W V_c = V_f \implies V_c(\mathcal{J}_i, :)^T \widetilde{w}_i = v_i \qquad \forall i \in \mathcal{F}$$

where  $\mathcal{J}_i$  is the set of column indices of the prescribed nonzeros of the *i*-th row of W. Note that this is another sequence of  $n_f$  small and dense, generally **underdetermined**, linear systems.



By collecting all the non-zero entries of W into the vector w and using Lagrange multipliers, the problem can be stated as the solution of the following saddle-point system:

$$\begin{bmatrix} K & B \\ B^T & 0 \end{bmatrix} \begin{bmatrix} w \\ \lambda \end{bmatrix} = \begin{bmatrix} f \\ g \end{bmatrix}$$

#### **Observations:**

- The size of K is typically much larger than the size of A (about 20 times larger in 3D mechanics);
- With a proper unknown ordering both K and B are block diagonal matrices;
- If *w* is numbered following the columns of *P* then *K* is block diagonal;
- If *w* is numbered following the rows of *P* then *B* is block diagonal;
- Unfortunately, they cannot be block diagonal at the same time;

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As in the original paper, we solve the minimization problem through a **restricted conjugate gradient**, which is a nullspace method. Define the projection orthogonal to B,  $\Pi_B = I - B(B^T B)^{-1} B^T$ , so that:

$$B^T y = 0 \qquad \forall y = \Pi_B v$$

Then, starting from a tentative prolongation  $W_0$  already satisfying the constraint,  $B^T w_0 = g$ , we apply the **conjugate gradient** method to the system:

$$\Pi_B K \Pi_B \Delta w = \Pi_B (f - K w_0)$$

to find a correction  $\Delta W$  to be applied to  $W_0$ :

$$W = W_0 + \Delta W$$

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During CG iteration, it is possible to monitor the decrease in energy and stop when a plateu is reached.

Some issues related to the proposed solution methd:

- The energy minimization is just a piece of the AMG set-up, so we cannot spend to much effort in it ⇒ it must converge quickly!
- Memory occupation is a big issue. K and B require a significant amount of memory and sometimes it is unavoidable to proceed in matrix-free mode;
- Preconditioning is crucial:
  - projected Jacobi,  $\Pi_B D_K \Pi_B$ , is enough for many problems;
  - incomplete Cholesky with very low fill-in converges faster but it is more difficult to apply, if K cannot be stored;

## Numerical results



#### Linear elasticity Benchmark:

Regular discretization of an elastic cube with tetrahedral FE.

| Matrix    | # of rows | # of non-zeroes |
|-----------|-----------|-----------------|
| C_4820    | 14,460    | 556,686         |
| C_35199   | 105,597   | 4,079,357       |
| C_246389  | 739,167   | 29,610,351      |
| C_1772489 | 5,317,443 | 222,268,213     |

We compare the following prolongation strategies:

- **EXTI** Extended+i (hypre);
- BAMG least squares fit;
- **SBAMG** BAMG + smoothing;
- **EMIN** energy minimization;





| Matrix    | P-type | $C_{gd}/C_{op}$ | $n_{it}$ |
|-----------|--------|-----------------|----------|
| C_4820    | EXTI   | 1.062/1.200     | 91       |
|           | BAMG   | 1.062/1.162     | 25       |
|           | SBAMG  | 1.062/1.404     | 17       |
|           | EMIN   | 1.062/1.383     | 12       |
| C_35199   | EXTI   | 1.068/1.254     | 209      |
|           | BAMG   | 1.067/1.211     | 65       |
|           | SBAMG  | 1.066/1.576     | 33       |
|           | EMIN   | 1.066/1.540     | 15       |
| C_246389  | EXTI   | 1.066/1.214     | 510      |
|           | BAMG   | 1.064/1.196     | 146      |
|           | SBAMG  | 1.064/1.608     | 55       |
|           | EMIN   | 1.064/1.517     | 22       |
| C_1772481 | EXTI   | 1.064/1.188     | 1286     |
|           | BAMG   | 1.062/1.189     | 270      |
|           | SBAMG  | 1.062/1.590     | 78       |
|           | EMIN   | 1.062/1.460     | 33       |



Figure: # of PCG iterations vs. grid size

## Numerical results





Figure: Relative residual,  $\frac{r_k}{r_0}$ , vs. iteration count, k

## Numerical results



| Matrix           | # of rows | # of non-zeroes |
|------------------|-----------|-----------------|
| Pflow_742k       | 742,793   | 37,138,461      |
| <i>Gear_167k</i> | 167,460   | 9,861,437       |
| S4A_700k         | 631,007   | 26,148,363      |
| Mech_447k        | 447,703   | 18,243,793      |



Figure: # of PCG iterations.



Figure: Total solution time [s].

## Conclusions



- Energy minimizing prolongation is a well-known concept in AMG, which, however has little application due to its cost and its difficult implementation;
- Our main contribution was to make energy minimization feasible by means of preconditioning and an effective implementation;
- Our algorithm has proved very effective on linear systems arising from mechanical problems and also in challenging problems arising from different applications.
- Though requiring a slightly more expensive set-up, this additional cost is largerly offset by a fast convergence.

Further work:

- Further improve preconditioning of restricted CG;
- Include energy minimization in the HPC AMG solver Chronos.

G. Isotton, M. Frigo, N. Spiezia and C. Janna, *Chronos: A general purpose classical AMG solver for high performance computing*, **SIAM Journal on Scientific Computing 43**, C335-C357, 2021.

## Thanks for your attention. Questions?