

# A study on adaptive algorithms for numerical quadrature on heterogeneous multicore and GPU based systems

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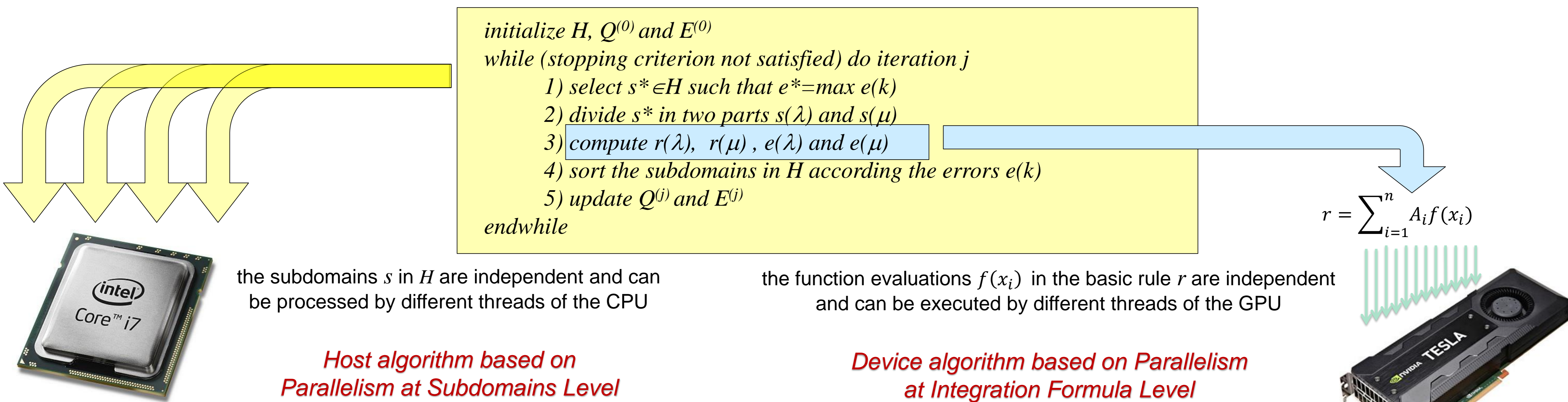
A parallel adaptive algorithm for the computation of a multidimensional integral on heterogeneous systems is described. Two different strategies have been combined together in a single algorithm: a first procedure is in charge of the load balancing among the threads on the multicore CPU and a second one is in charge of an efficient execution on the GPU of the computational kernel. Experimental results on a system with a quad-core CPUs Intel Core I7 950 @ 3Ghz and two GPUs NVIDIA C1060 have been achieved.

**PROBLEM:**  $I(f) = \int_U f(t_1, \dots, t_d) dt_1 \dots dt_d$   
 $U = [a_1, b_1] \times \dots \times [a_d, b_d] \quad 2 \leq d \leq 10$

**Basic rule:**  $r = \sum_{i=1}^n A_i f(x_i) \sim \int_U f(t_1, \dots, t_d) dt_1 \dots dt_d$   
 $O(10^3) \leq n \leq O(10^4)$  (Genz & Malik rule)

## Adaptive algorithm:

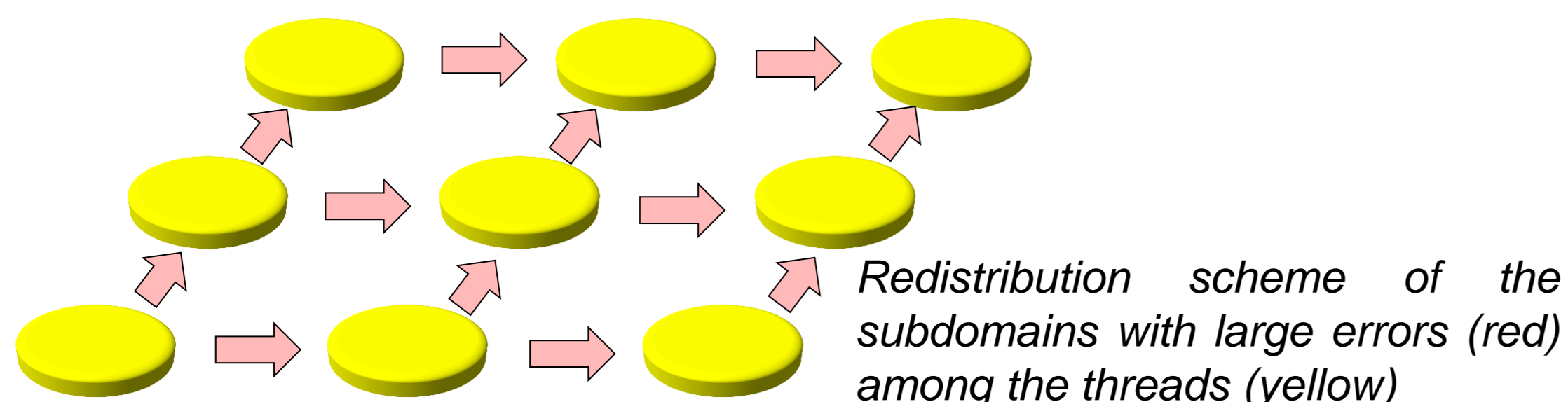
An iterative procedure that refines the integration domain  $U$ , evaluating the quadrature rule  $r$  only in the subdomains  $s^*$  with the largest error  $e^*$  in a suitable data structure  $H$



## Main issue: design of a scalable H

The access to a single centralized  $H$  produces fast numerical convergence but several synchronizations among all threads (therefore very poor scalability)

**Solution:** No centralized data structure! The threads are logically organized according to a 2-dimensional periodical mesh  $M_2$ . Each thread  $T_i$  manages a private sub-structure  $H_i$ . Then it attempts to share its item  $s^*_i \in H_i$  with largest error  $e^*_i$ , only with the neighbor threads in the mesh  $M_2$ .

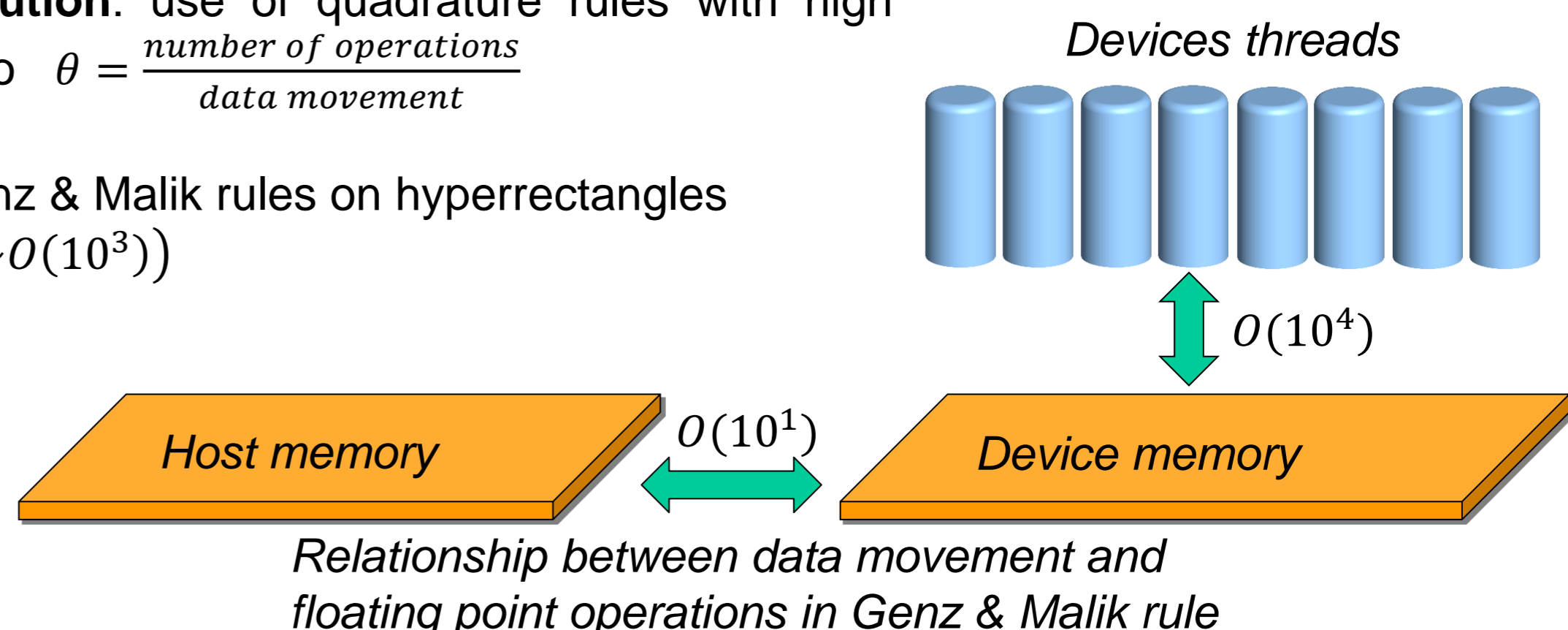


## Main issues: management of the fine grain parallelism

1) Summation is a hard to optimize kernel (e.g. many idle threads)  
**Solution:** use of the reduce functions in CUDA toolkit 4.0

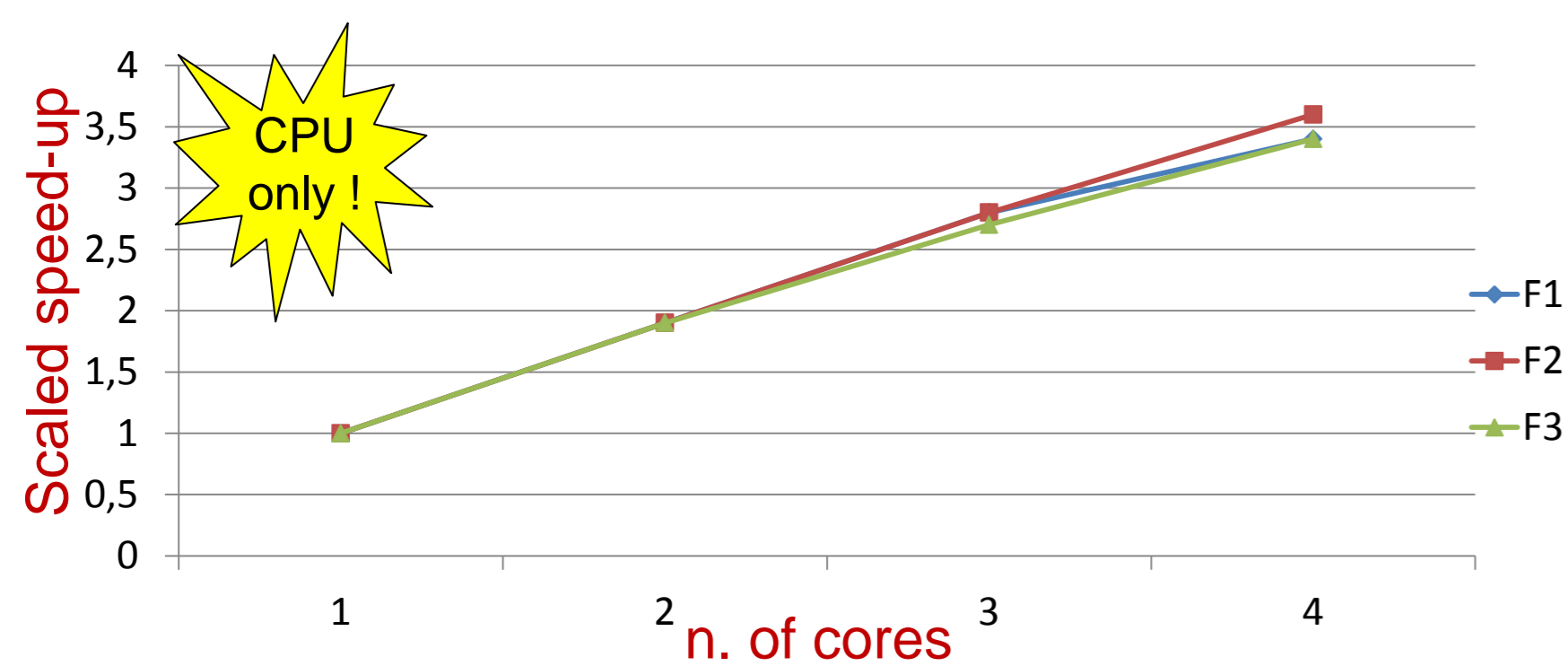
2) Low bandwidth between device and host memories  
**Solution:** use of quadrature rules with high ratio  $\theta = \frac{\text{number of operations}}{\text{data movement}}$

Genz & Malik rules on hyperrectangles ( $\theta \sim O(10^3)$ )

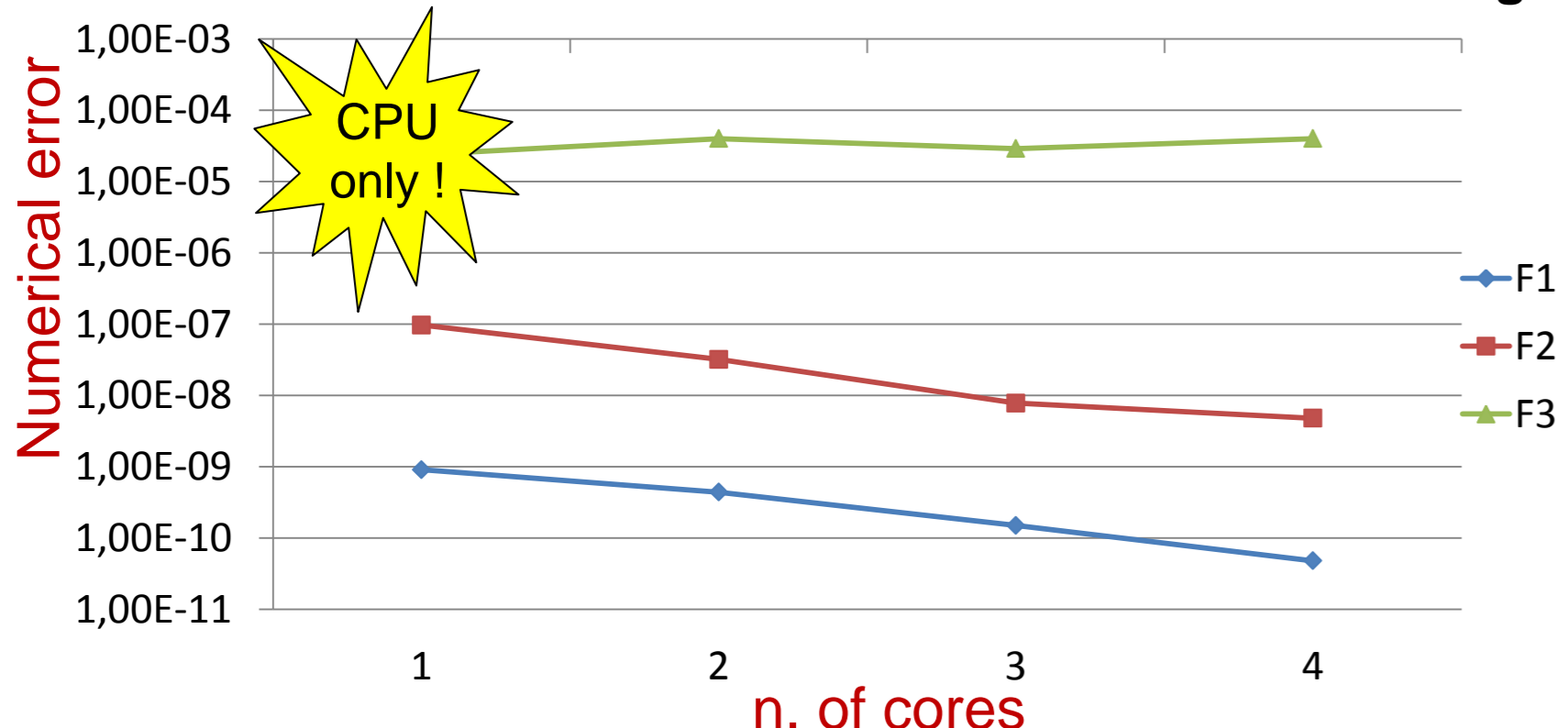


## TEST RESULTS

### Scalability test (fixed time) on the CPU when the number of cores grows

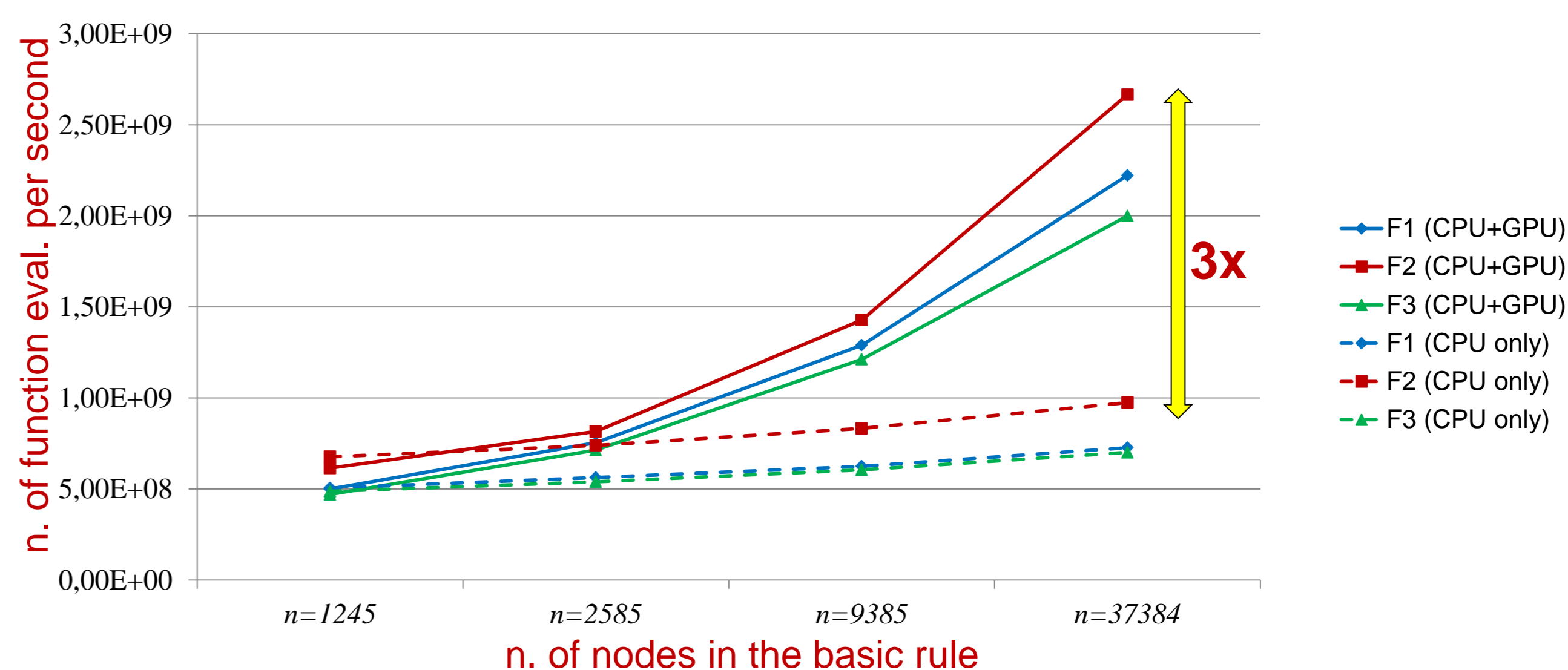


### Error reduction on the CPU when the number of cores grows



10 functions in each family  
 $F1 = \cos(2\pi\beta_1 + \sum_{i=1}^d \alpha_i x_i)$  oscillating functions  
 $F2 = (1 + \sum_{i=1}^d \alpha_i x_i)^{-d-1}$  corner peak functions  
 $F3 = \exp(-\sum_{i=1}^d \alpha_i |x_i - \beta_i|)$   $C^0$  functions  
 • Basic rule with  $n=1245$  function evaluations in  $d=10$  dimensions  
 •  $10 \times 10^6$  function evaluations per core (4016 iterations)

### Performance gain of the algorithm with (solid line) and without (dashed line) the use of GPU



## CONCLUSIONS

- We presented a heterogeneous multicore CPU/GPU algorithm with a performance gain of 3x with respect to a traditional quadrature adaptive algorithm running just on current homogeneous multicore CPUs.
- The approach demonstrates the utility of graphics accelerators for multidimensional quadrature mainly with a large number of function evaluations of the basic rule ( $n > 10^4$ ) and in large dimension ( $d > 10$ ).
- Our approach can be combined with other levels of parallelism (e.g. cluster level)



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